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Simulation Theorems in Multi-valued Modal μ -Calculus

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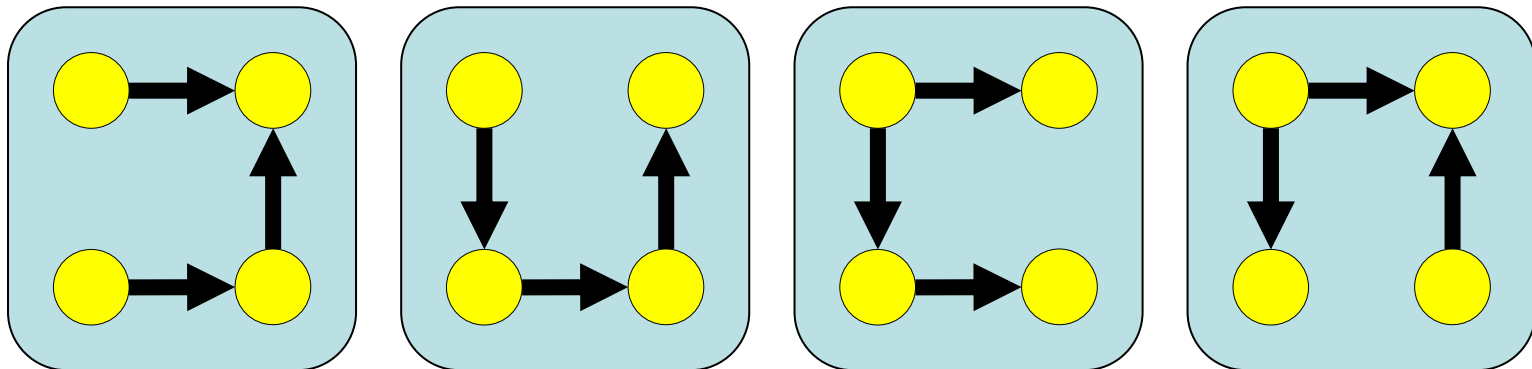
Yoshiki Kinoshita (AIST/CVS)

4th VERITE March 6, 2007

Motivation

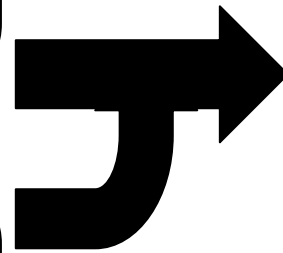
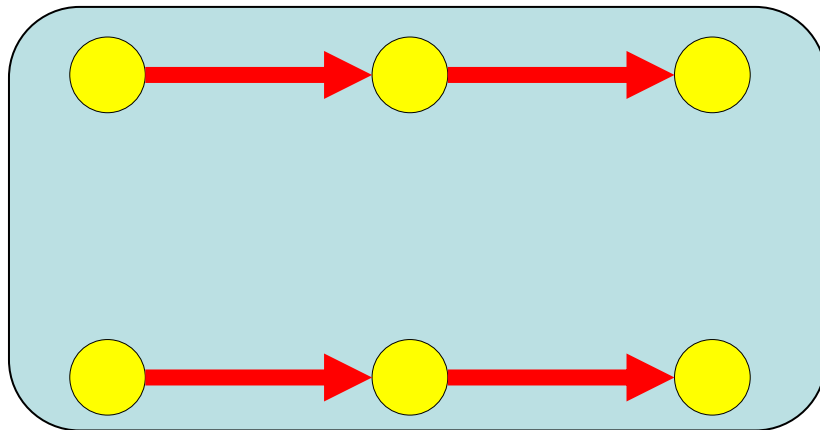
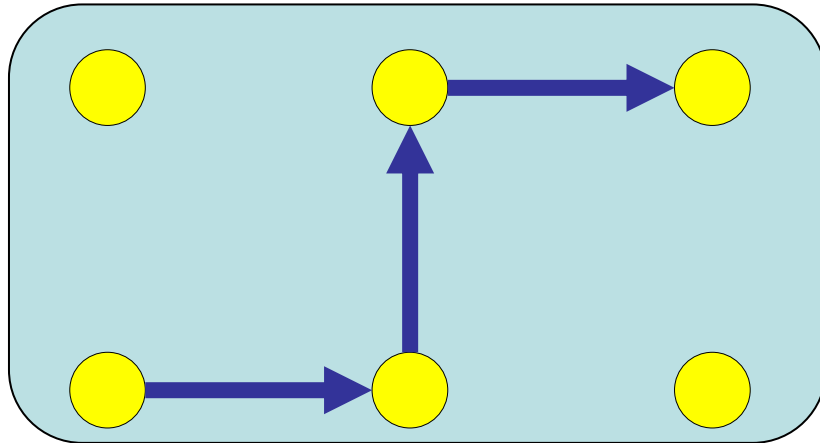
“Refinement of Models” in Model Checking

- Model Checking = Modeling + Checking
- Tatsumi and Kameyama tried to get minimal one among models checked successfully.
- They needed a number of model checking.

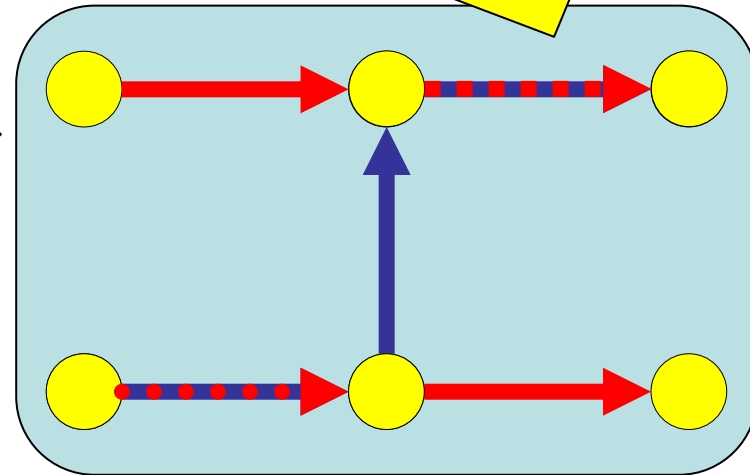


They wanted to perform a number of model checking all at once.

Superposition of Models



Let L be the set of weights of transitions



From $2 = \{T, F\}$ to general L

- Transition System, Kripke Model, Simulation
- State semantics of Modal μ -Calculus, Simulation Theorem
 - De Morgan algebra [Tatsumi-Kameyama 2006]
 - Complete Heyting algebra [This talk]
- Path semantics of Linear Modal μ -calculus, Simulation Theorem
 - Complete Heyting algebra + condition [This talk]

Why complete Heyting algebra ?

- Sets and binary relations form a **category**.
- **L must** be a complete Heyting algebra for sets and binary L-valued relations to form a category [Johnstone 2002].

Complete Heyting algebra

is $(L, \leq, \vee, \wedge, \Rightarrow)$ satisfying the following.

1. (L, \leq) is a **partially ordered set**.
2. An arbitrary subset of L has the **join** (so, also the **meet**).
3. $a \wedge b \leq c \Leftrightarrow b \leq a \Rightarrow c$

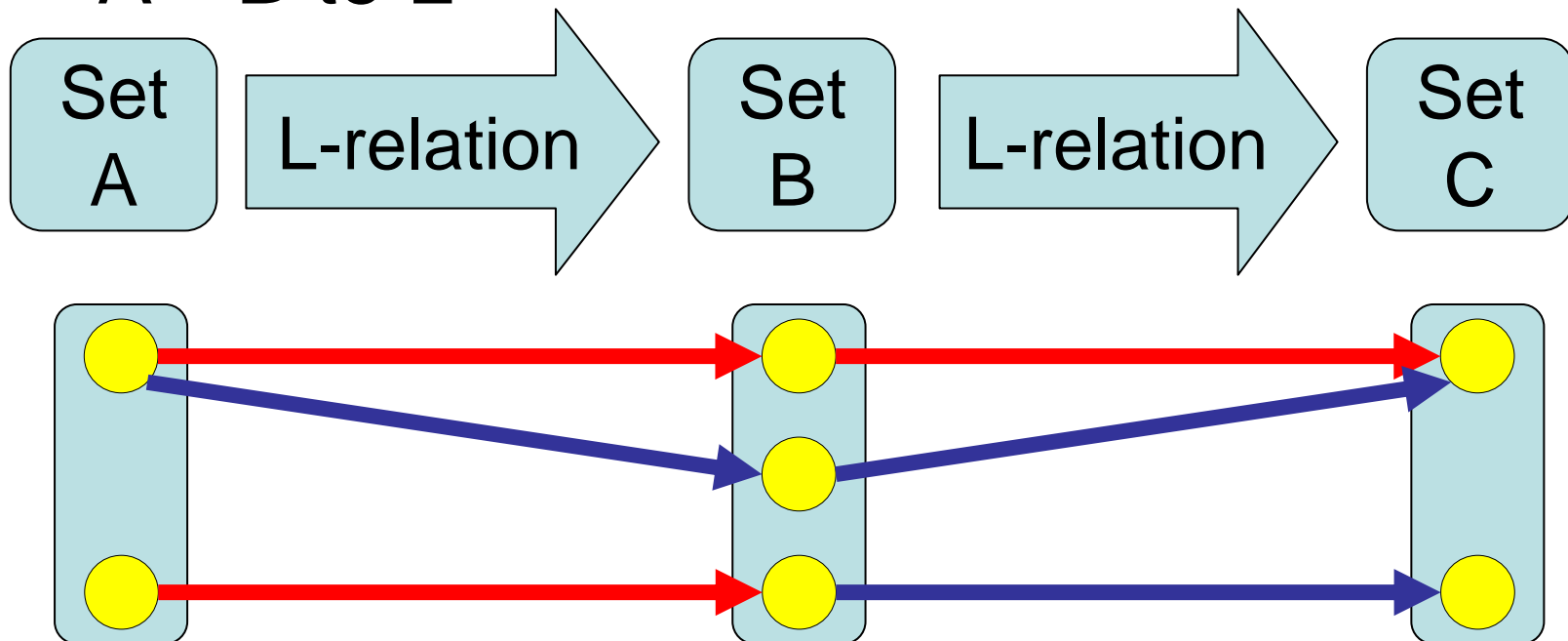
Example:

$2, 2 \times 2, \dots, 2^n, \dots$

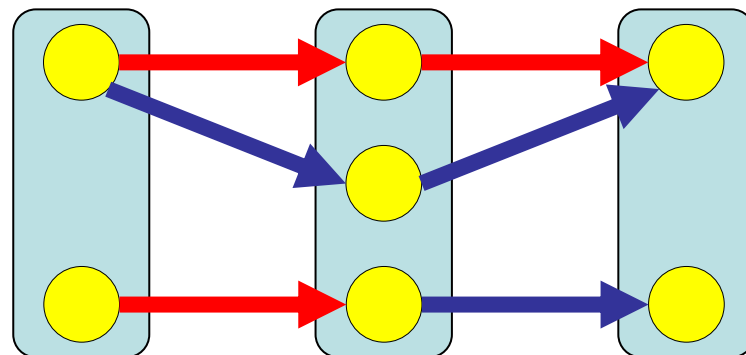
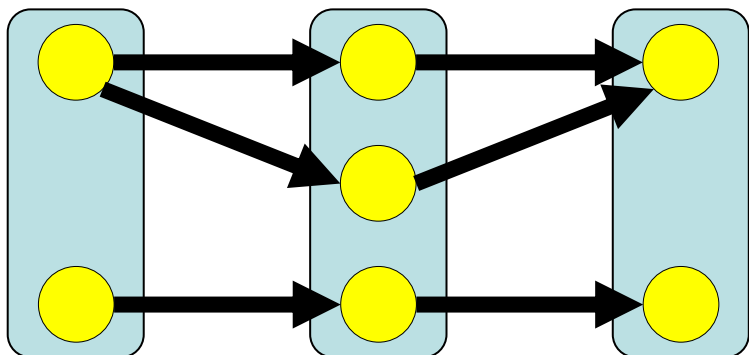
The open sets of a topological space

Category of L-valued relations

- Objects are sets
- Arrows from A to B are functions from $A \times B$ to L

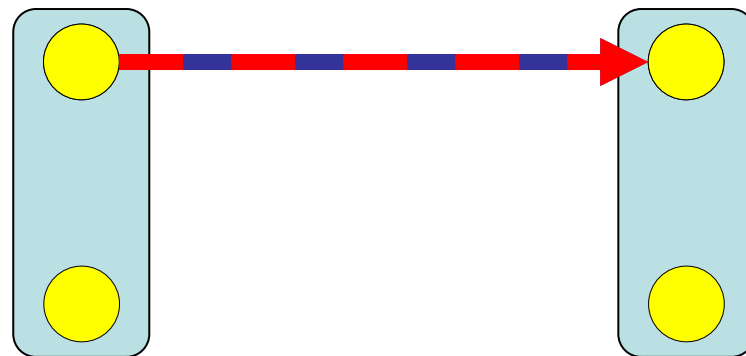
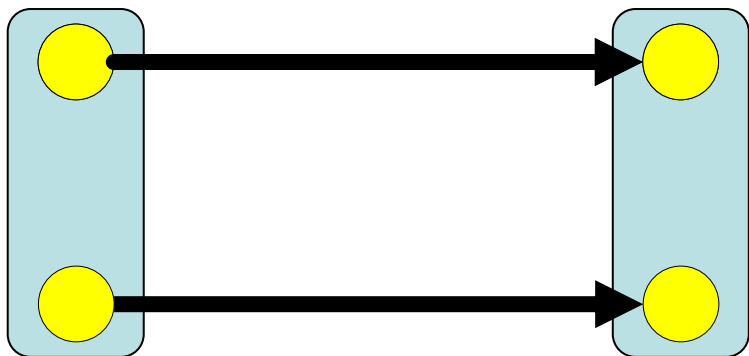


Composition: $L=2$ and $L=2 \times 2$

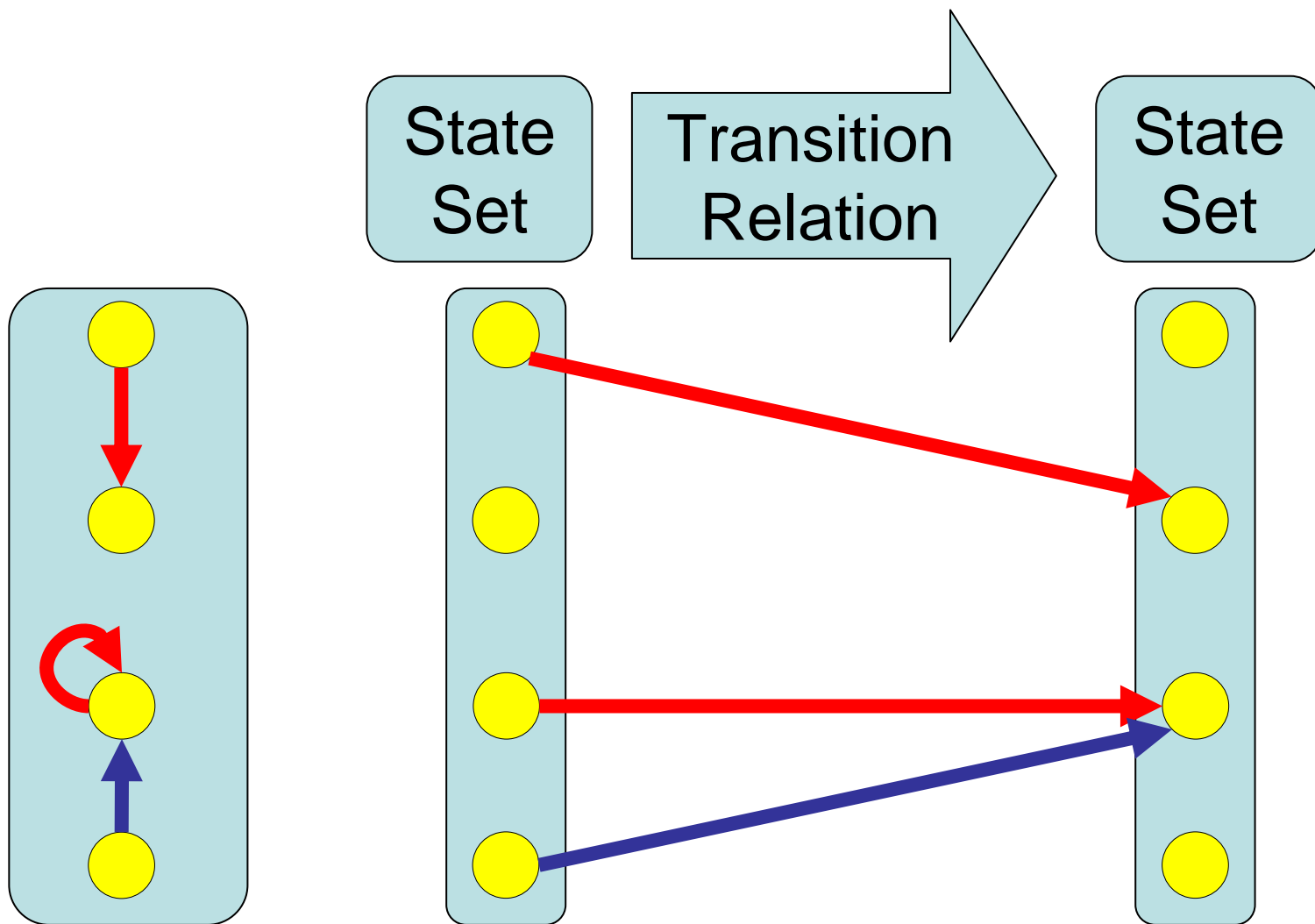


Compose

Compose

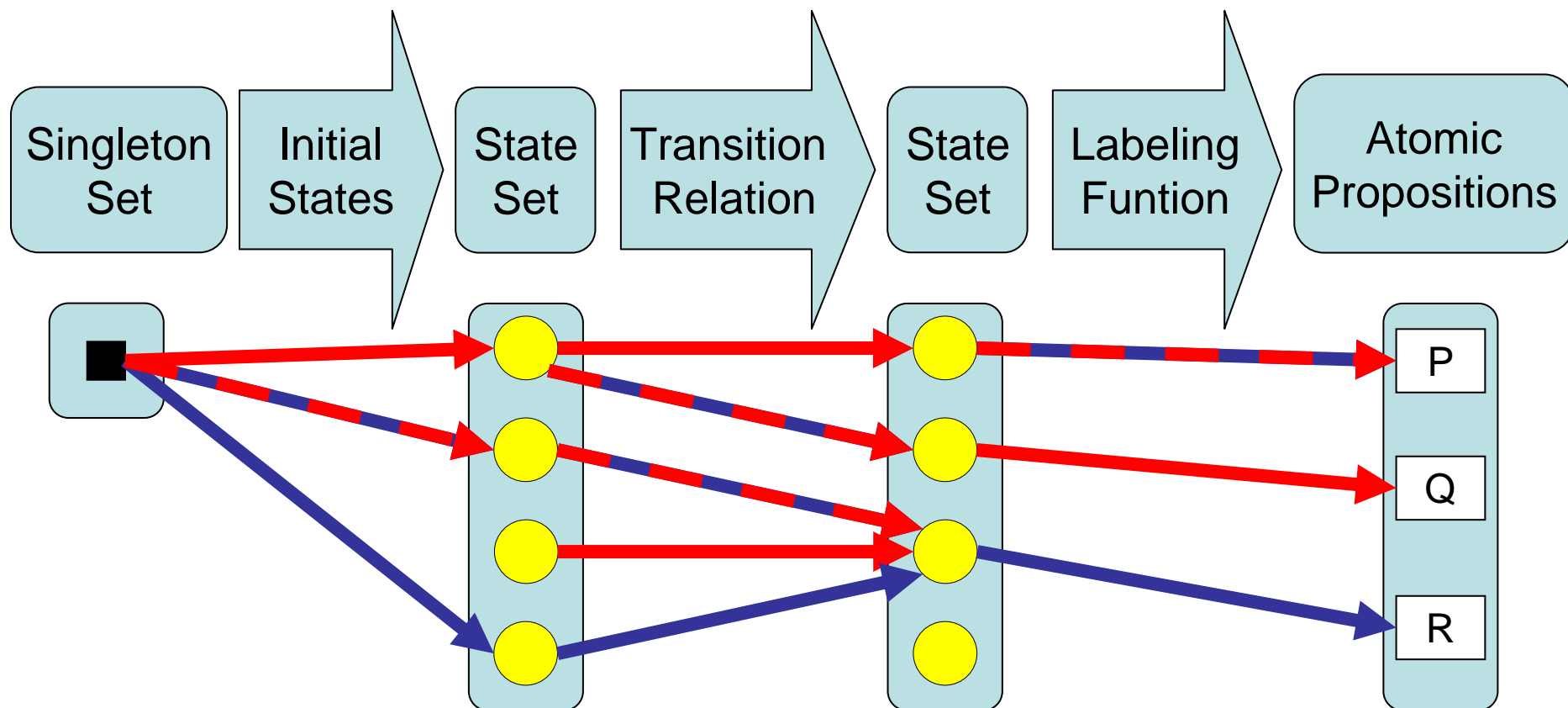


L-valued Transition System

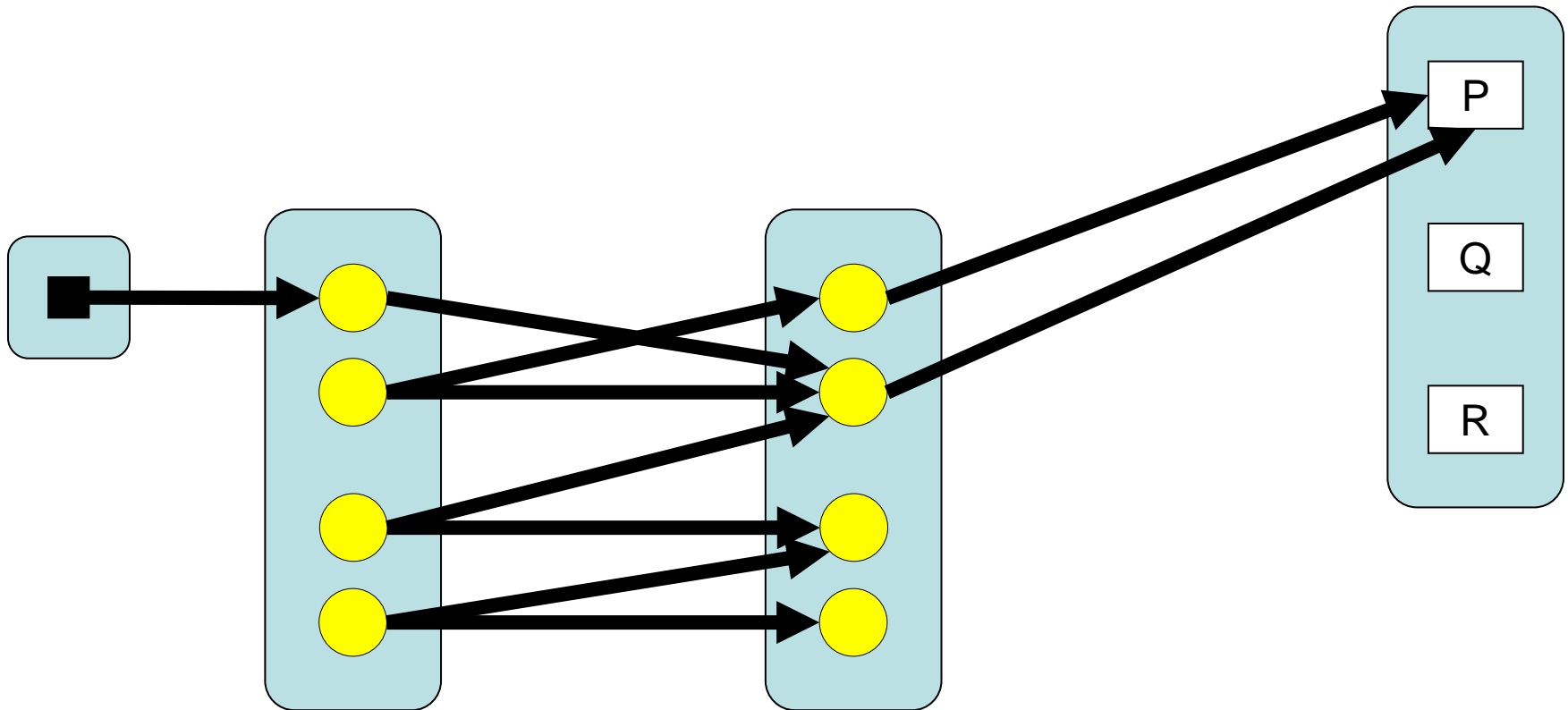


L-valued Kripke model

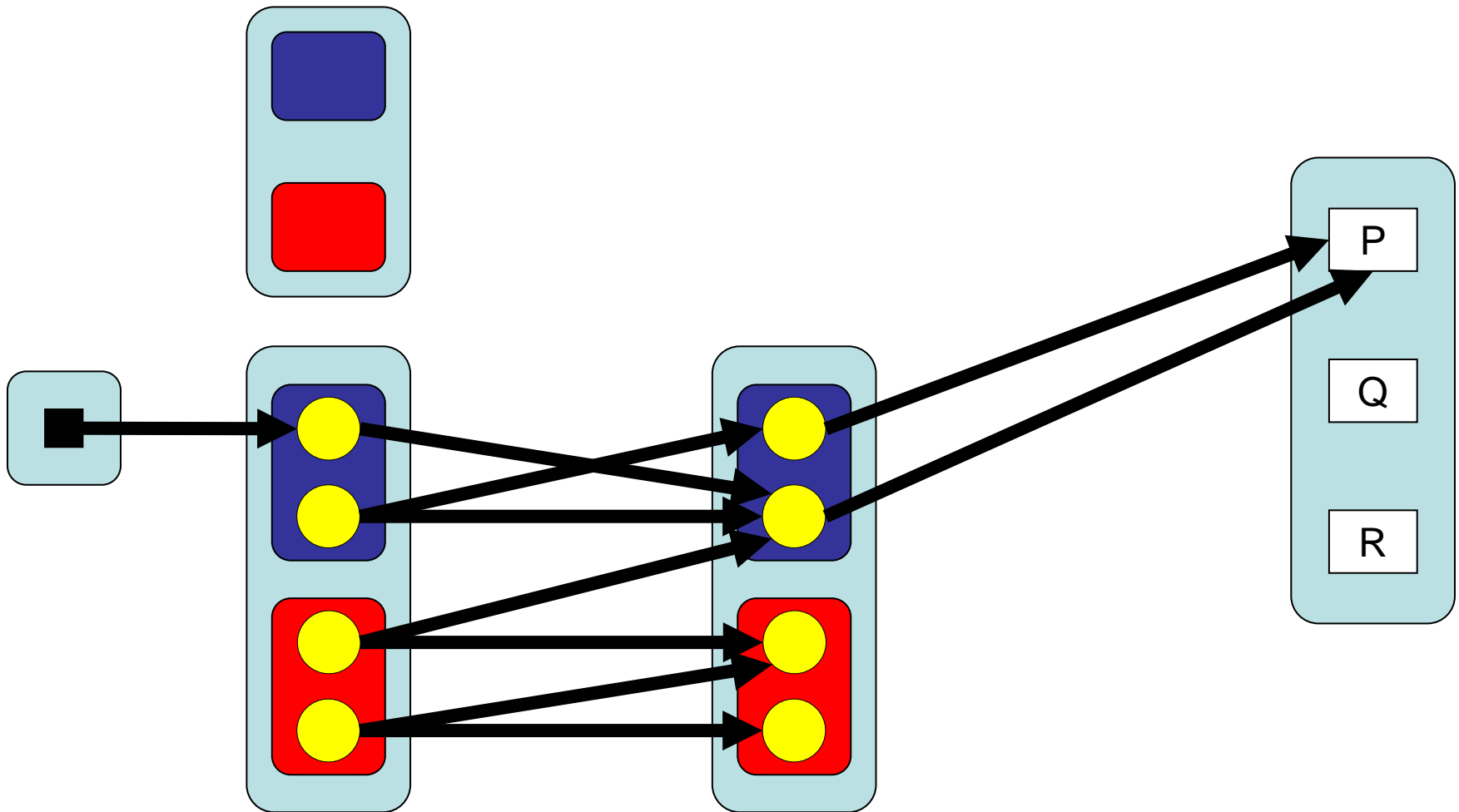
consists of the following **L-relations**.



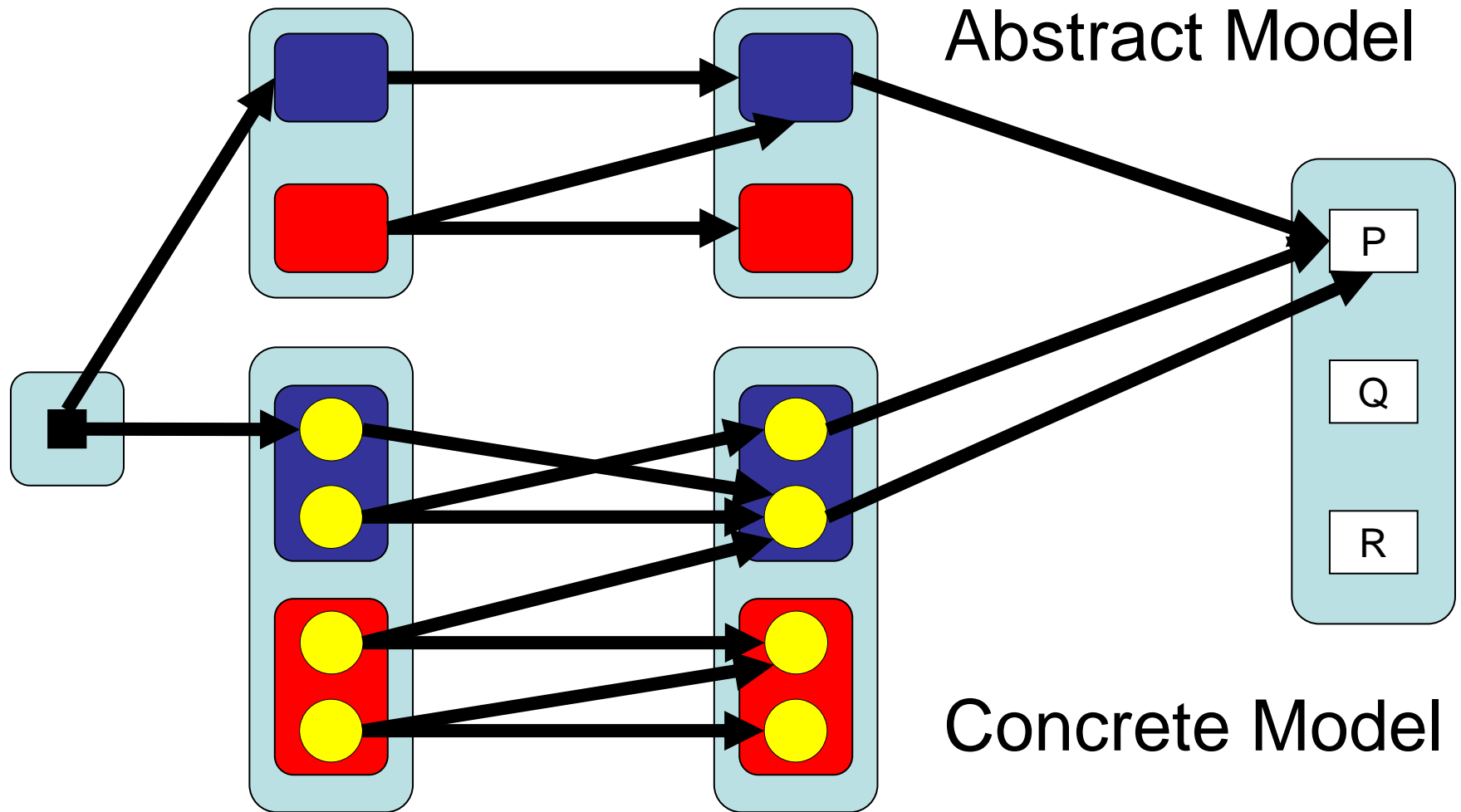
2-valued Simulation



2-valued Simulation

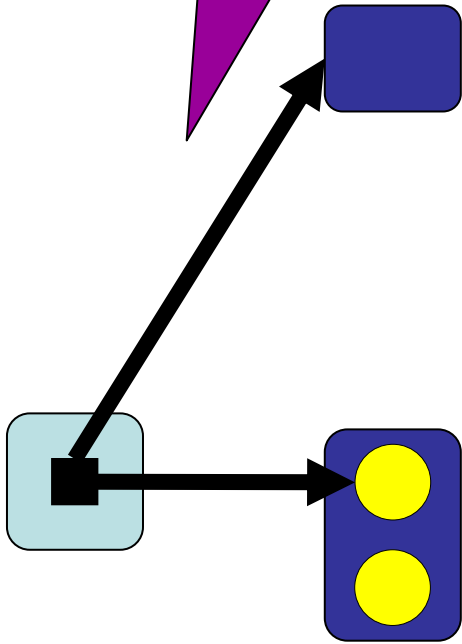


2-valued Simulation



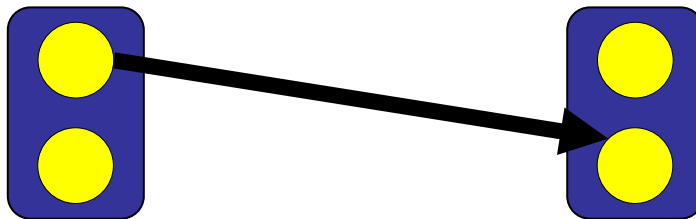
Initial
States

2-valued Simulation

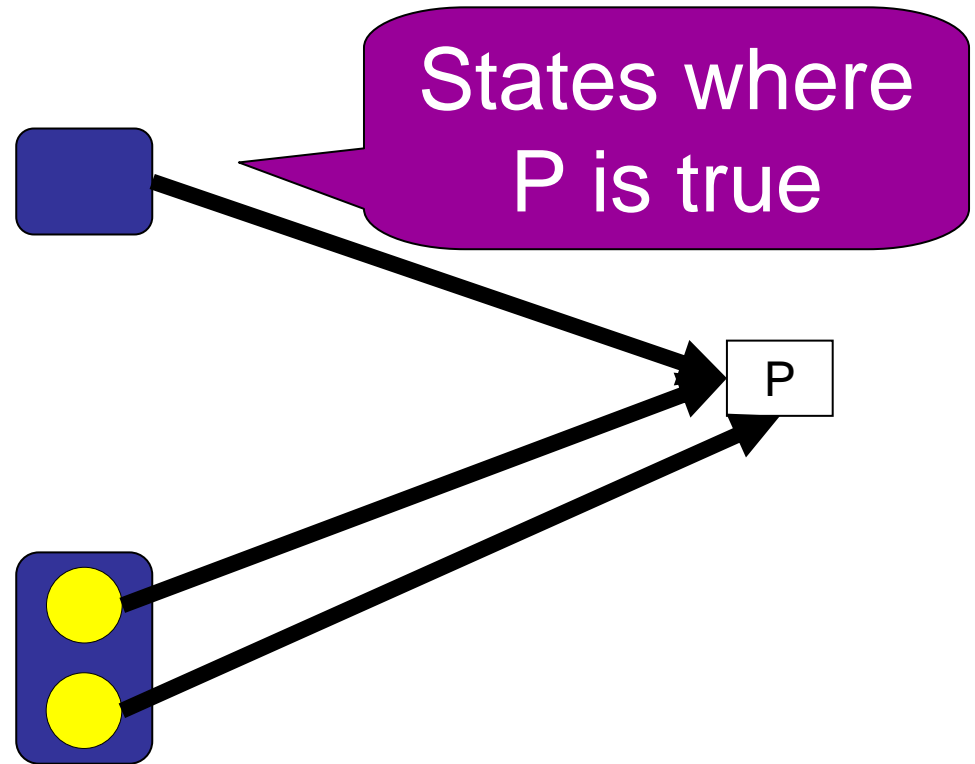


2-valued Simulation

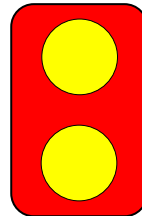
Transition



2-valued Simulation



2-valued Simulation



States where
P is false

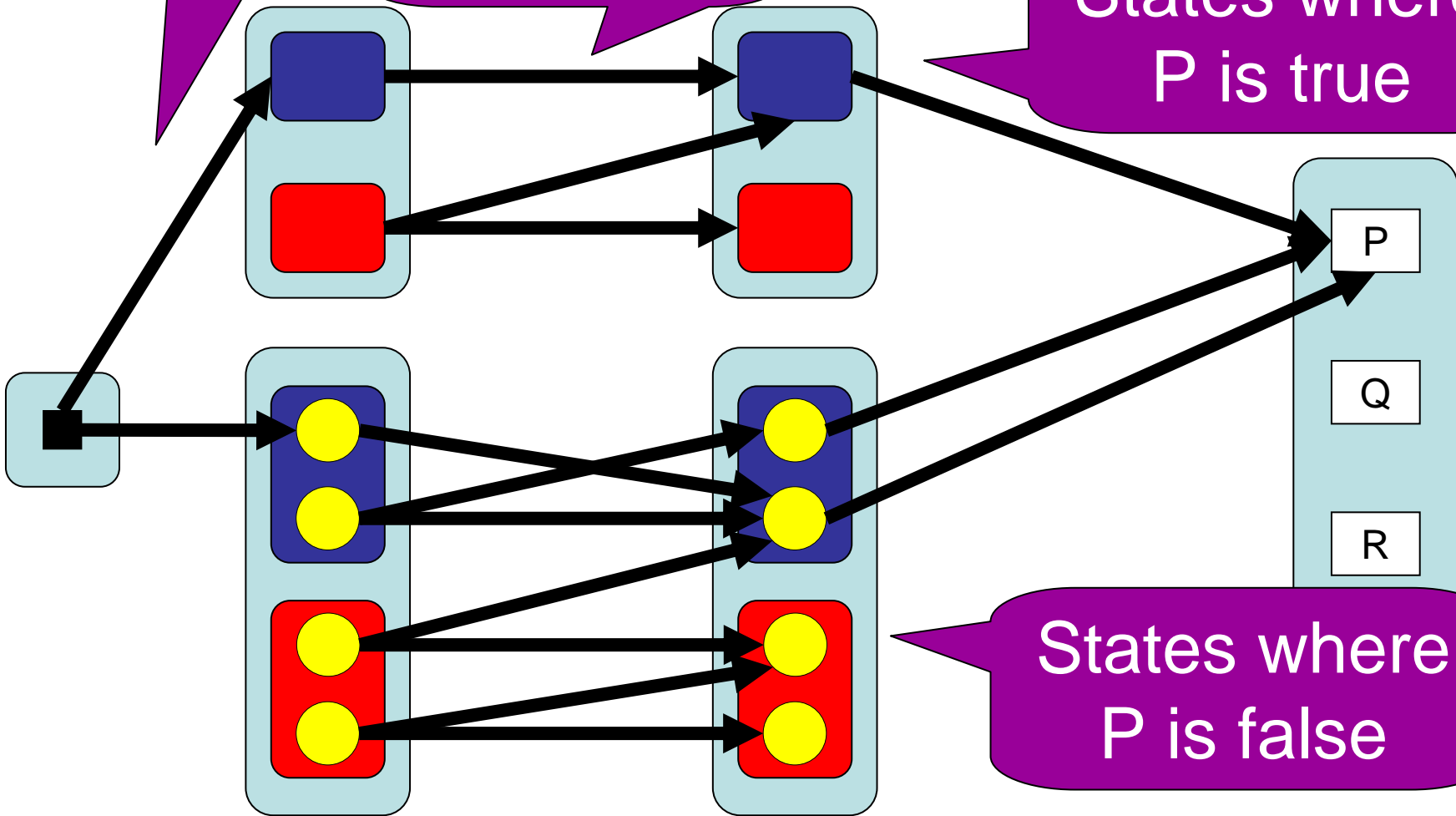
2-valued Simulation

Initial States

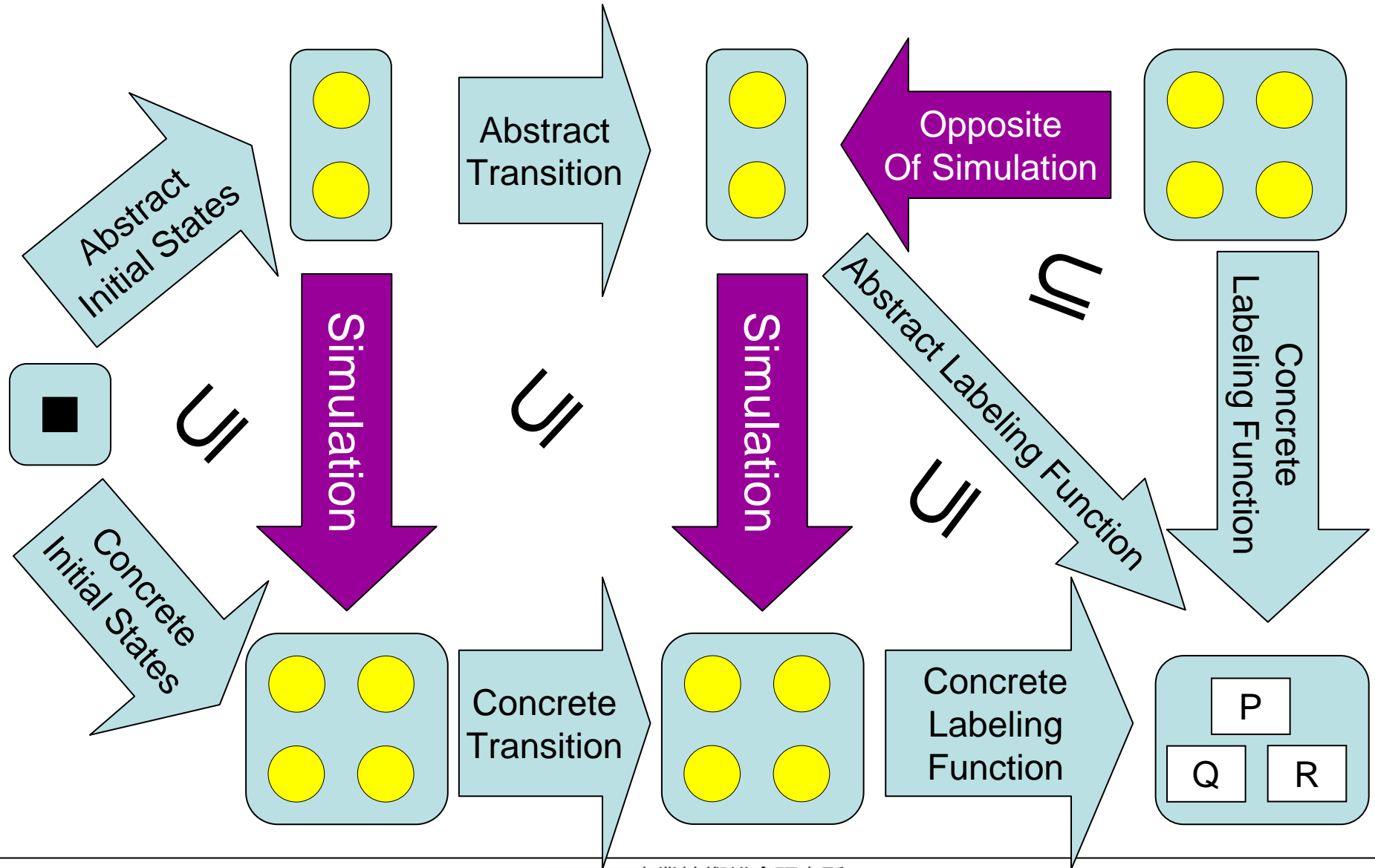
Transition

States where P is true

States where P is false



L-valued Simulation



L-valued State Semantics

Modal μ -Calculus

$$\psi ::= p \mid \perp \mid \top \mid \psi \vee \psi \mid \psi \wedge \psi \mid \psi \Rightarrow \psi \\ \mid x \mid \mu x. \psi \mid \nu x. \psi \mid \diamond \psi \mid \square \psi$$

$K, s, V \models \psi$ is **an element of L**

- Natural definition (no details in this talk)
- Intuitionistic version

$$K, s, V \models \psi \quad \neq \quad K, s, V \models (\psi \Rightarrow \perp) \Rightarrow \perp$$

Simulation Theorem

For any simulation,
if the abstract model satisfies ψ ,
then the concrete model satisfies ψ .

- When ψ has no \square in the negative positions and no \diamond in the positive positions
- Example: $\nu X.P \wedge \square X$
“P always globally holds”.

This theorem holds in L-valued context.

L-valued Path Semantics

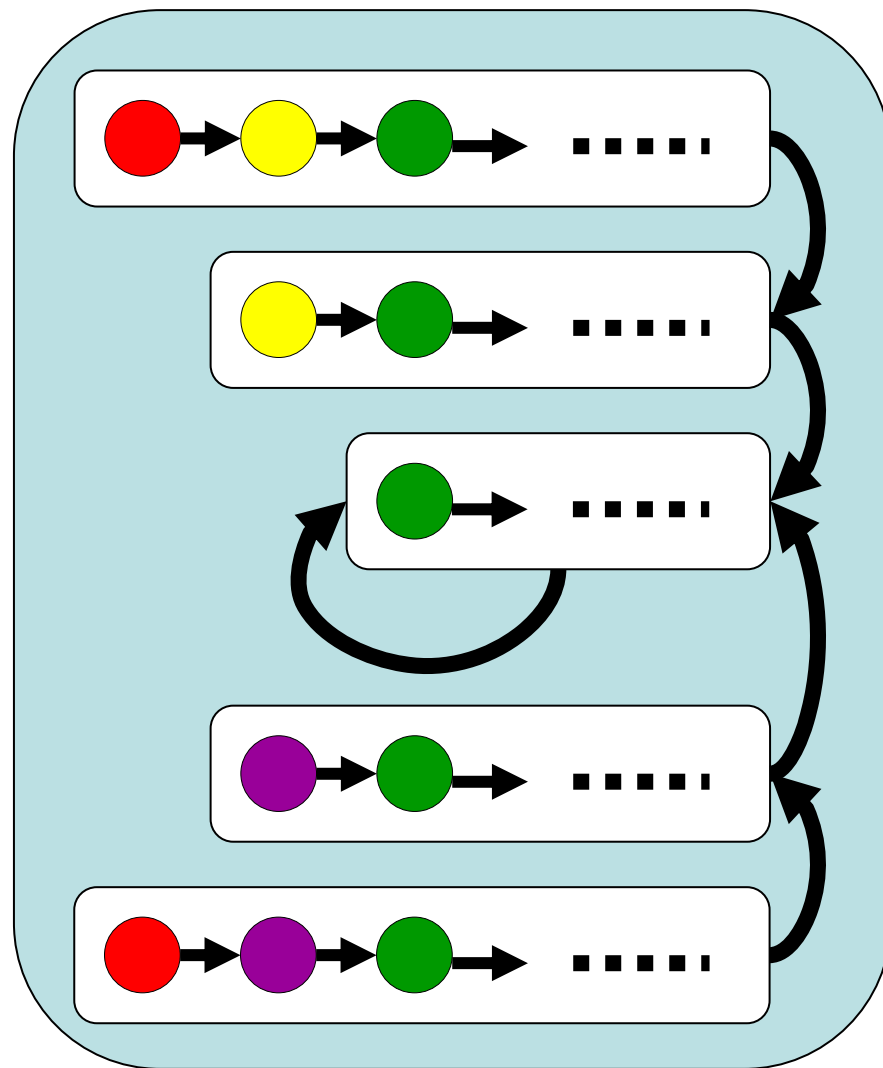
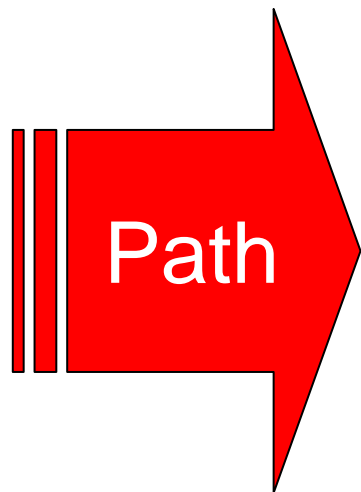
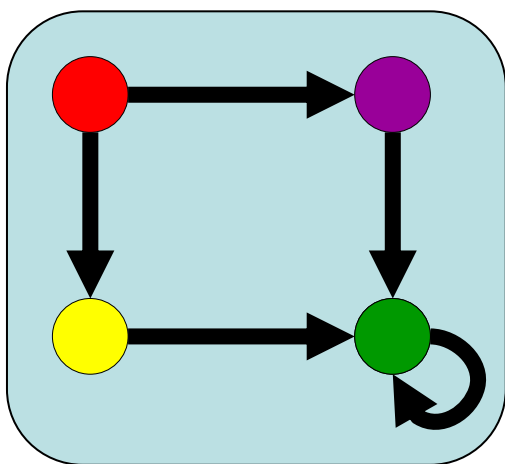
Linear Modal μ -Calculus (generalization of LTL)

$$\psi ::= p \mid \perp \mid \top \mid \psi \vee \psi \mid \psi \wedge \psi \mid \psi \Rightarrow \psi$$
$$\mid x \mid \mu x. \psi \mid \nu x. \psi \mid \text{Next } \psi$$

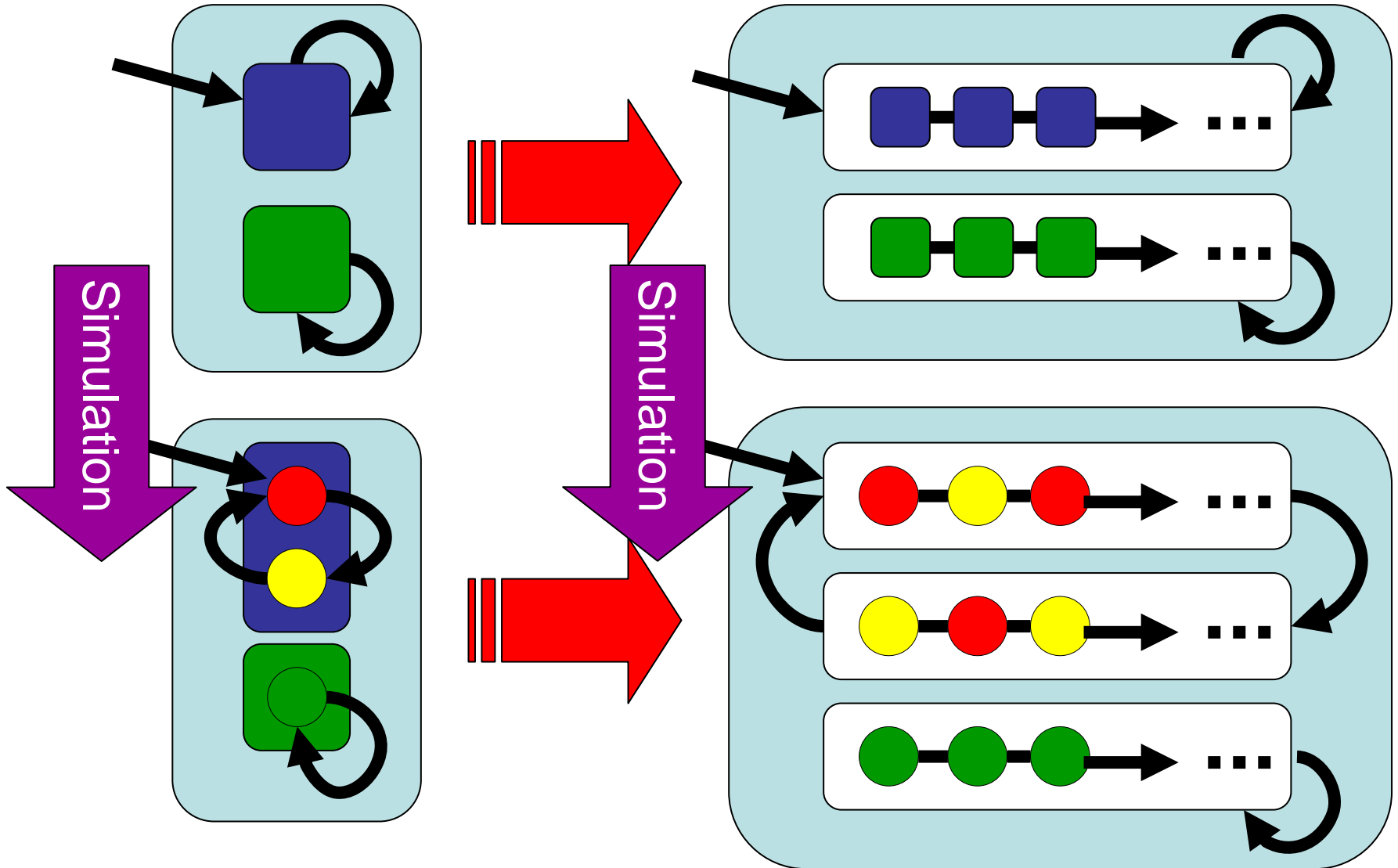
$K, \pi, V \models \psi$ is defined for a path π .

2-valued Path Semantics

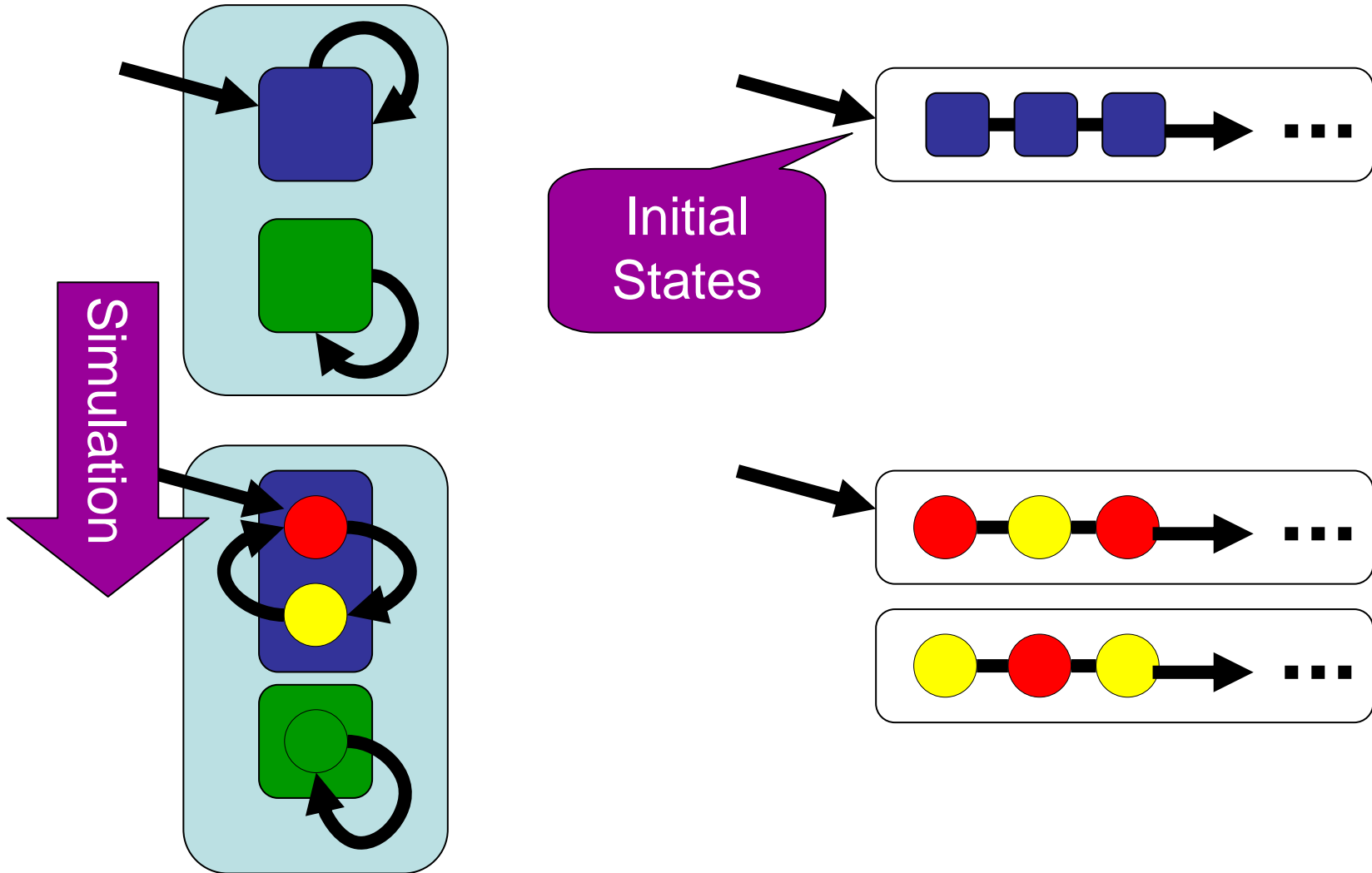
Path Semantics =
Path Construction
+ State Semantics



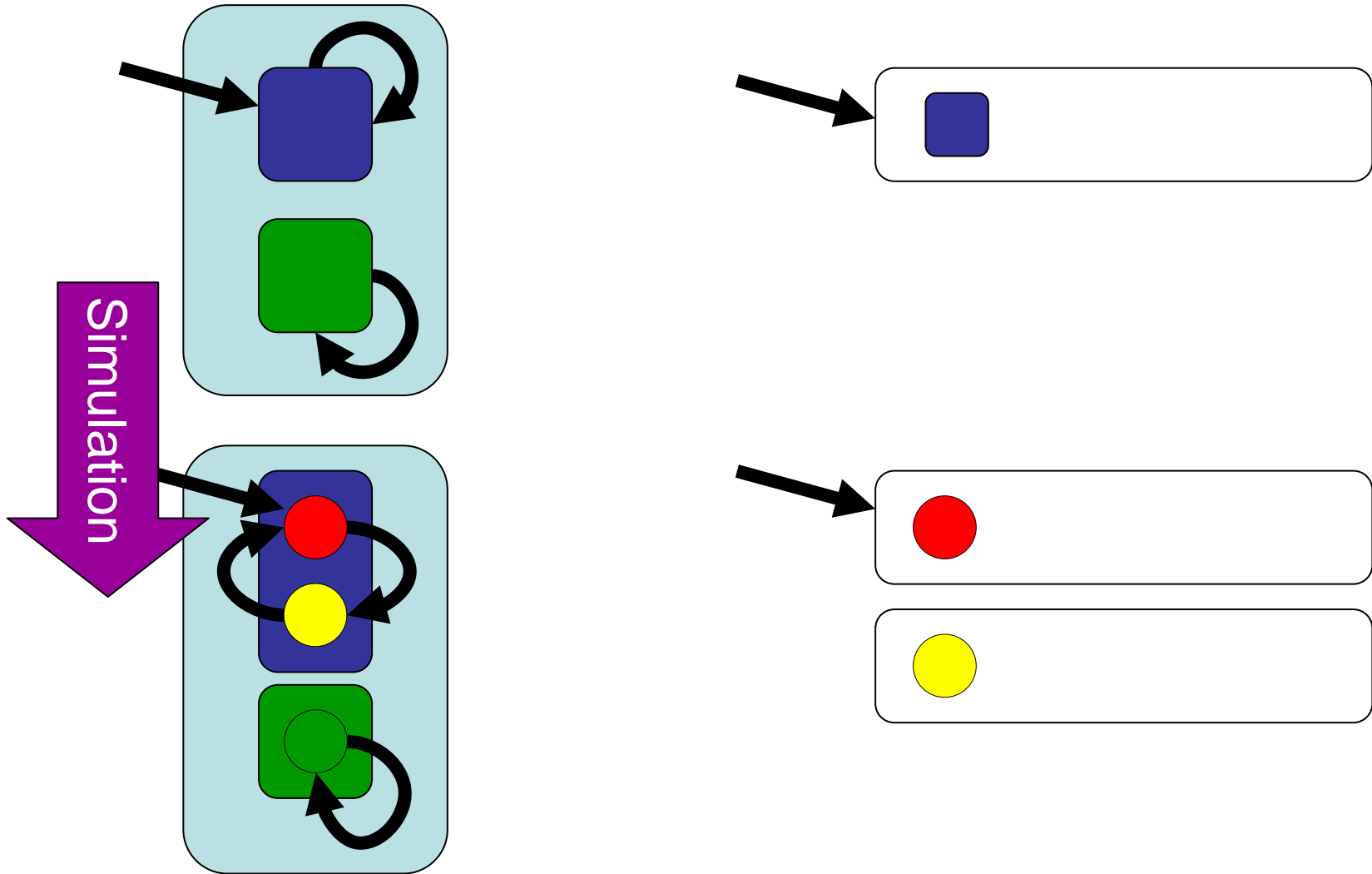
Simulation is lifted



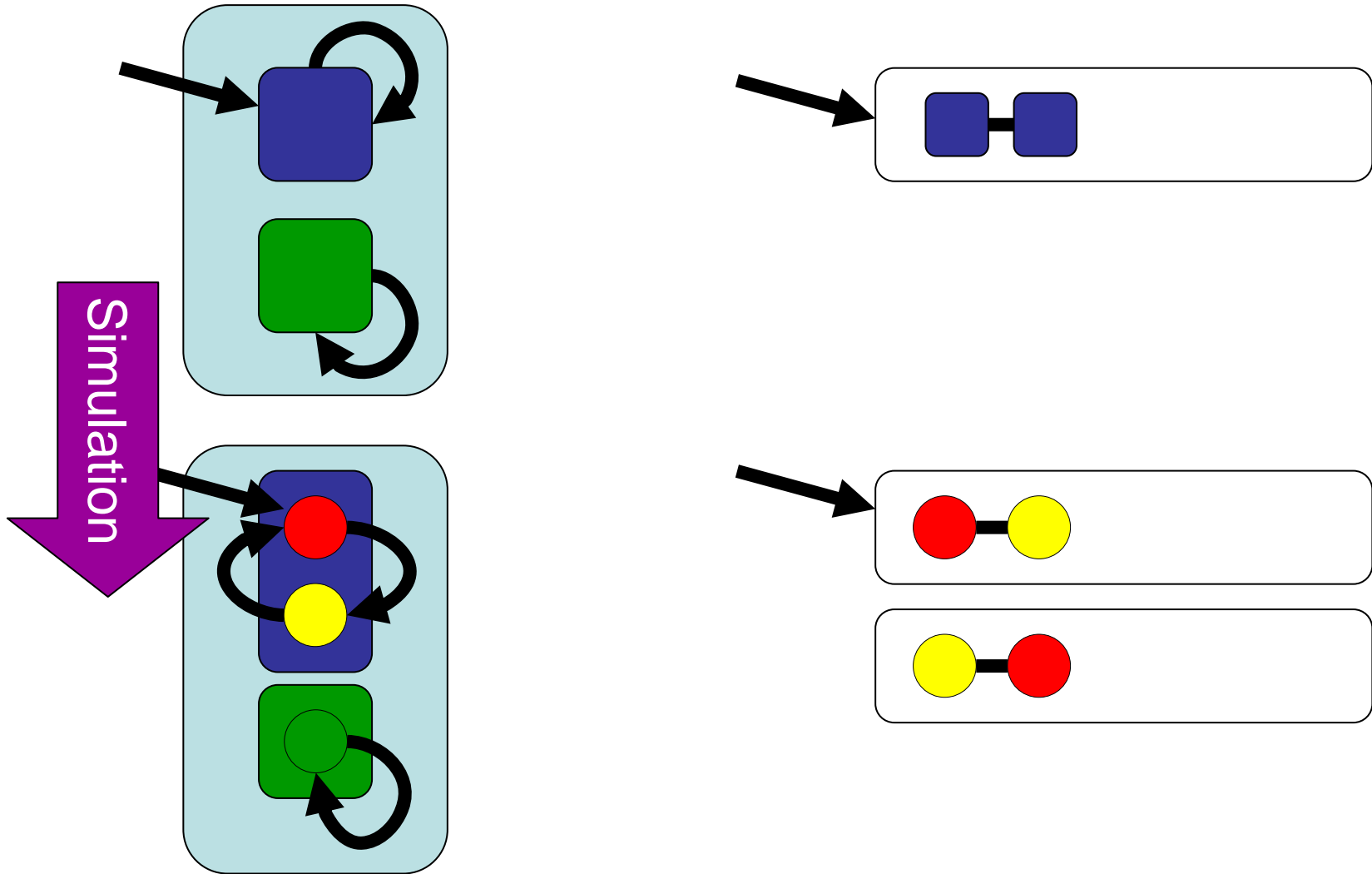
Simulation is lifted



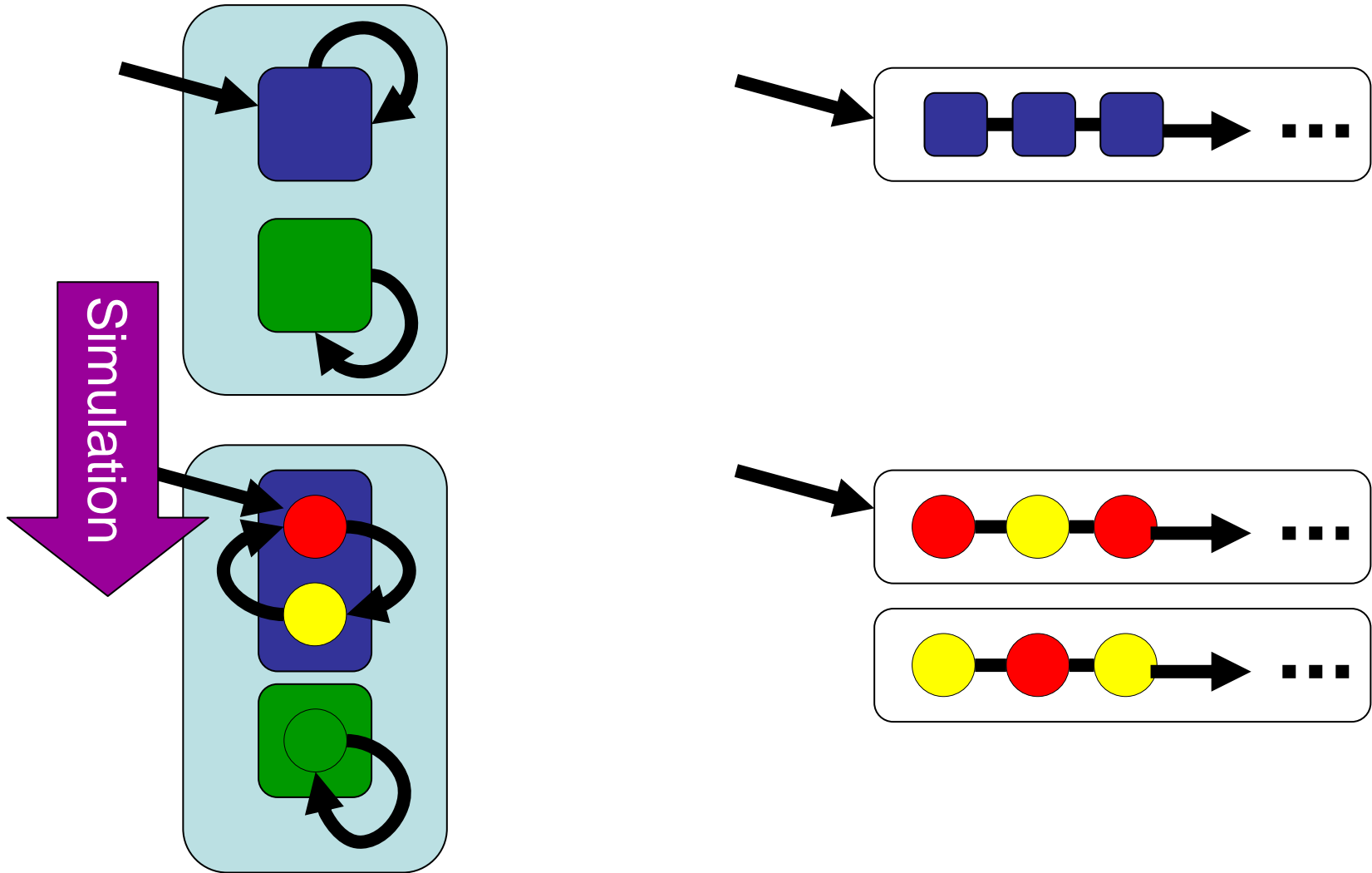
Simulation is lifted



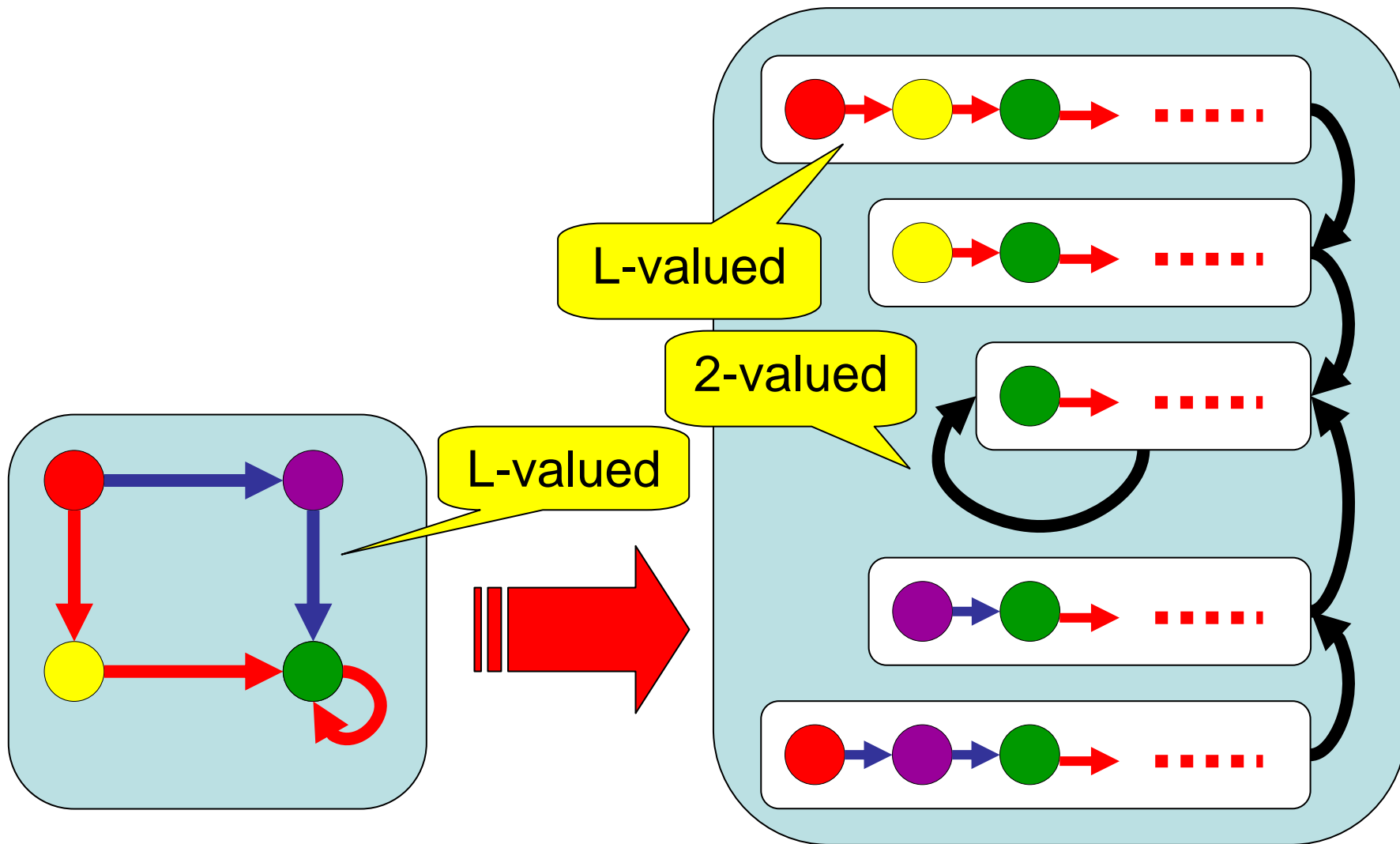
Simulation is lifted



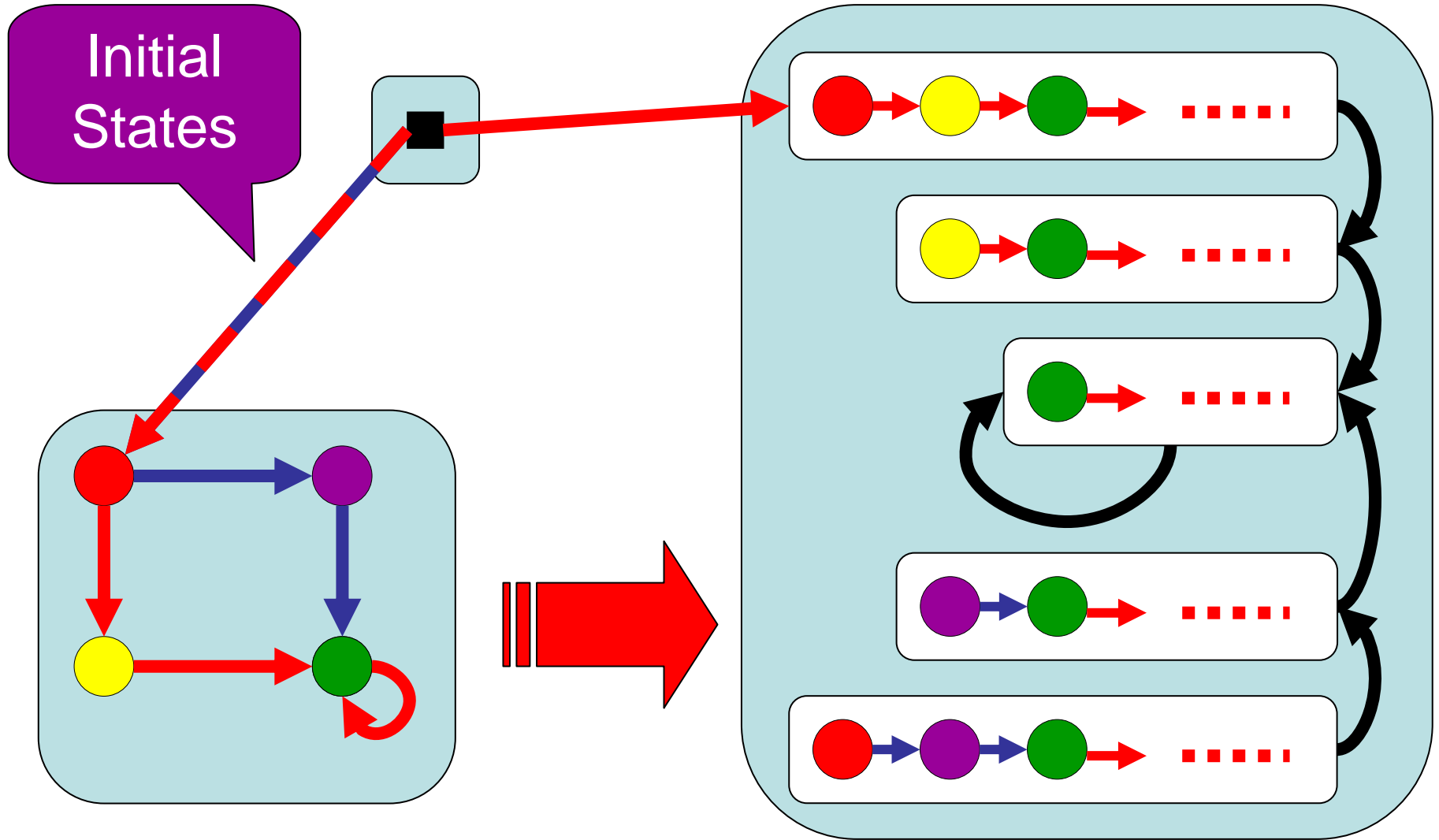
Simulation is lifted



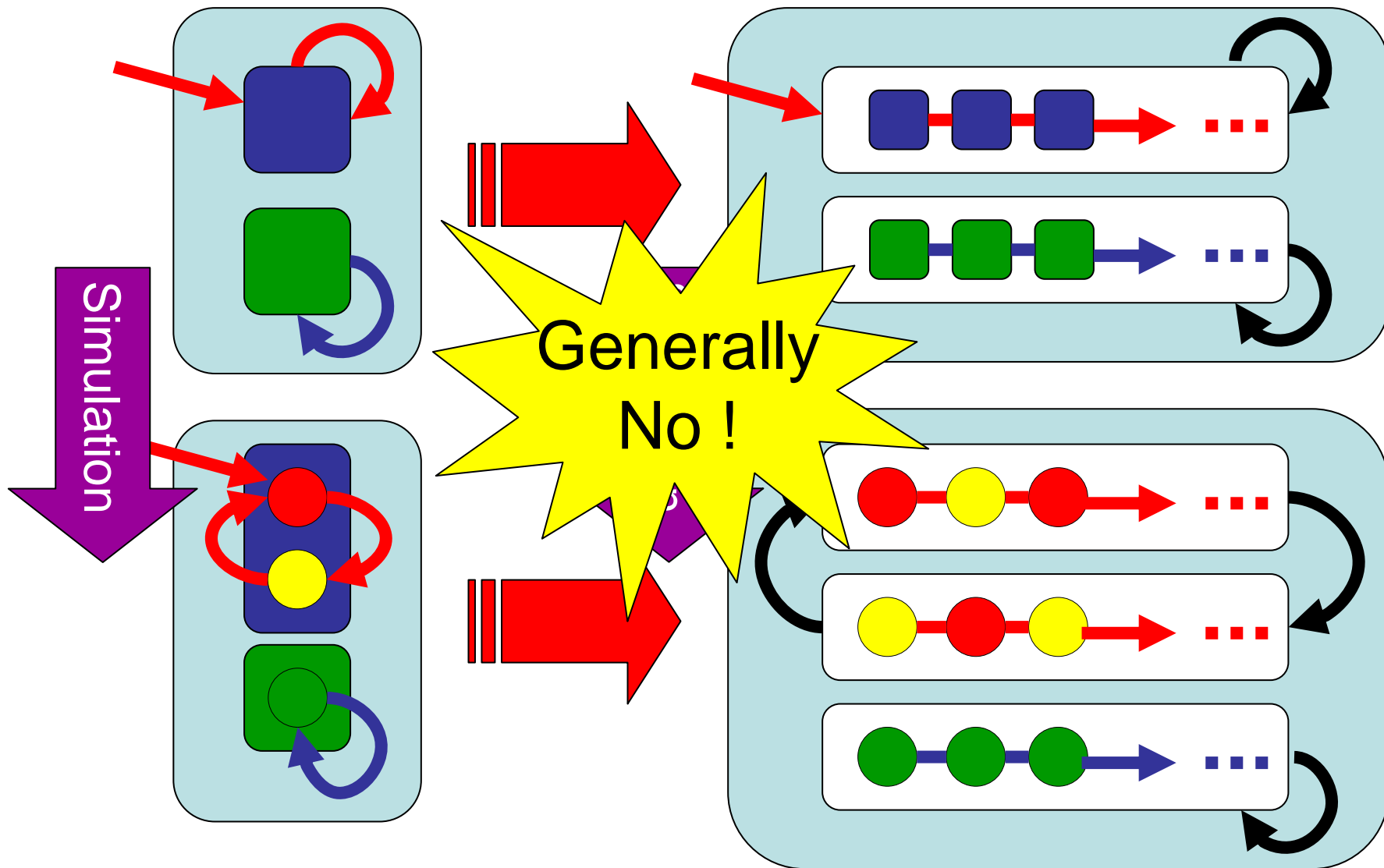
L-valued Path Semantics



L-valued Path Semantics



Is simulation lifted ?



Is simulation lifted ?

- **No.** We found a counterexample.
- We gave a sufficient condition:
 - A simulation is lifted if L is **the open sets** of a topological space and **closed for countable intersections**.
 - Examples: power sets, $\text{Nat} \cup \{\omega\}$
- Under the condition, the simulation theorem for path semantics holds.

Conclusion

- Complete Heyting algebra valued
 - Transition System, Kripke Model, Simulation
 - State Semantics for Modal μ -Calculus, Simulation Theorem
- Under our new condition
 - Path Semantics for Linear Modal μ -Calculus, Simulation Theorem

Future Work

- To relate this work to fuzzy relations or probabilistic relations