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Author(s)	Zhao, Dan; Dauch, Axel; Matsumoto, Tad
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Modulation Doping for Repetition Coded BICM-ID with Irregular Degree Allocation

Dan Zhao

Japan Advanced Institute of Science
and Technology, Japan
Email: dan.zhao@jaist.ac.jp

Axel Dauch

Information Technology
University of Ulm, Germany
Email: Axel.Dauch@gmx.de

Tad Matsumoto

Japan Advanced Institute of Science and
Technology, Japan
Email: matumoto@jaist.ac.jp
and
Center for Wireless Communication at
University of Oulu, Finland
Email: tadashi.matsumoto@ee.oulu.fi

Abstract—A new BICM-ID technique that makes efficient use of a mixture of standard Gray and anti-Gray extended mapping, combined with irregular degree-allocated repetition code with check nodes. It is shown that with the proposed technique, very low rate codes that can achieve turbo cliff at a very low E_b/N_0 value range can well be designed, and that the EXIT chart analysis and chain simulation results are exactly consistent with each other.

Index terms - extended mapping; repetition code; irregular node allocation; EXIT analysis; turbo cliff; capacity;

I. INTRODUCTION

In Bit-Interleaved Coded Modulation with Iterative Detection/Decoding (BICM-ID) [1], transmitter is comprised of a concatenation of encoder and bit-to-symbol mapper separated by a bit interleaver, and at the receiver, iterative processing is invoked, according to the standard turbo principle.

Performances of BICM-ID have to be evaluated by the convergence and asymptotic properties [2], which are represented by the threshold signal-to-noise power ratio (SNR) and bit error rate (BER) floor, respectively. In principle, since BICM-ID is a serially concatenated system, analyzing its performances can rely on the area property [3] of the extrinsic information transfer (EXIT) chart. Therefore, the transmission link design based on BICM-ID falls into the issue of matching between the demapper and decoder EXIT curves. Various efforts have been made in the seek of better matching between the two curves for minimizing the gap while still keeping the tunnel open, aiming, without requiring heavy detection/decoding complexity, at achieving lower threshold SNR and BER floor. Reference [4] introduces a technique that makes good matching between the detector and decoder EXIT curves using low density parity check (LDPC) code in multiple input multiple output (MIMO) spatial multiplexing systems.

It has been believed that a combination of Gray mapping and Turbo or LDPC codes achieves the best performance comparing with other combinations for BICM-ID. However, Ref. [5] proposes another approach towards achieving good matching between the two curves by introducing different mapping rules, such as anti-Gray mapping, which allows the use of even simpler codes, such as convolutional codes, and with which still it is possible to achieve BER pinch-off (corresponding to the threshold SNR). Another good idea

that can provide us with design flexibility is extended mapping (EM) presented by [6], [7], [8] where with 2^m -QAM, more than m bits, l_{map} , are allocated to one signal point in the constellation. With this mapping rule, the left-most (LM) point of the demapper EXIT function has a lower value than with the Gray mapping, but the right-most (RM) point becomes higher. With this setting, the demapper EXIT function achieves better matching with even weaker codes such as small memory length convolution codes, and of the most practical importance is the fact that hardware complexity of the modulator and demodulator remains the same as that with the original 2^m -QAM.

In Ref [9], we proposed techniques that combine EM with an extremely simple code, repetition code, concatenated with single parity check code with irregular degree allocation. Even with this very simple combination, the EXIT curves of the demapper and decoder match well, and we can achieve BER pinch-off exactly at a designed SNR. The complexity for the Ref [9]'s proposed technique is extremely low, because EM does not require higher order modulation format and no iterations are needed in the decoder itself. Even with such simple structure, near-capacity performance can be achieved [9].

However, this technique is not suitable in designing BICM-ID with low rate codes that achieve BER pinch-off at very low SNR. This is because the LM point of the demapper EXIT function, corresponding to decoder feedback mutual information MI=0, is already very low, and there is not enough room to further lower the EXIT LM point while avoiding the intersection between the demapper and decoder curves. To solve this problem, this paper introduces modulation doping technique which mixes the labeling rules (=bit patterns allocated to each constellation point) for the extended anti-Gray mapping and the standard Gray mapping, at a certain ratio. Since the demapper EXIT function with Gray mapping is relatively flat, its LM point has the higher MI value than EM mapping rules. Therefore, with modulation doping the demapper EXIT LM point is lifted up from that without doping, and the amount depends on the doping ratio.

With the modulation doping technique combined with repetition coded EM BICM-ID with irregular degree allocation, we can achieve more degrees-of-freedom in choosing the parameter values so that we can achieve better matching between the demapper and decoder EXIT curves.

The paper is organized as follows. The system model is described in Section II. The basic structures of the technique presented by [9] are summarized in Section II; The structures introduced in [9] has three forms; (A) is comprised of only EM and repetition code, (B) comprised of variable and check nodes with regular degree allocation, and (C) comprised of EM BICM-ID with irregular degree allocation. The notations, (A), (B), and (C), are commonly used to describe Sub-section Indexes in this paper to identify the structures. Their convergence performances are evaluated in Section III using EXIT chart, and chain simulation results are shown in Section IV. Section V further extends the structures described in Section III to the case where modulation doing is combined. Results of the EXIT analysis and BER simulations are presented to verify the advantageous points of the combined use of EM and modulation doping. It is shown that the EXIT analysis and BER results are exactly consistent each other even with such low rate code cases where BER threshold happens at very low SNR range. Finally Section VI concludes this paper.

II. SYSTEM MODEL WITHOUT DOPING

A schematic diagram of the BICM-ID system is shown in Fig. 1.

Transmitter

The binary bit information sequence \underline{u} to be transmitted is encoded by (A) a simple repetition code, (B) a single parity check code where a single parity bit is added to every d_c-1 information bits, followed by a repetition code. d_c is referred to as check node degree. The structure (C) is the same as (B), but the repetition times d_v , referred to as variable node degree, take different values in one whole block (transmission frame); if d_v takes several different values in a block, such code is referred to as having irregular degree allocation. It is assumed that throughout the paper d_c takes only one identical value as in [4].

The coded bit sequence, encoded by the encoder (A), (B), or (C), is first bit-interleaved, segmented into l_{map} -bit segments, and then the each segment is mapped on to one of the 2^m constellation points for modulation. Since $l_{map} > m$ with the EM technique, more than one labels having different bit patterns in the segment are mapped on to each constellation point. However, there are many possible combinations of bit patterns to be allocated to the constellation points. This paper uses the labels assigned to the each constellation point, obtained by [7], so that, with full *a priori* information, the mutual information (MI) between the coded bit and the demapper output extrinsic LLR at the right most point of the demapper EXIT curve is maximized. Fig. 2 shows the optimal labeling for 4-QAM with $l_{map}=3$ and 5. The complex-valued signal modulated according to the mapping rule, referred to as Ψ , is finally transmitted to the wireless channel.

Channel

This paper assumes frequency flat block fading additive white Gaussian noise (AWGN) channel. If the channel exhibits frequency selectivity due to the multipath propagation, the receiver needs an equalizer to eliminate the

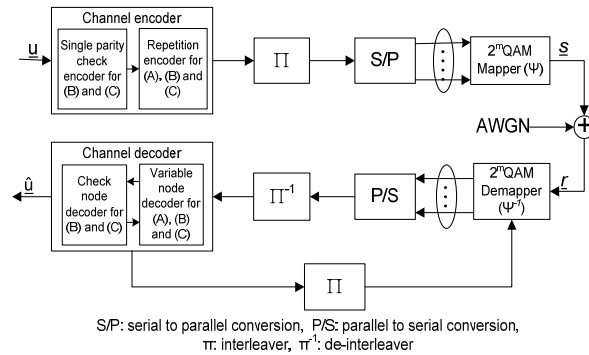


Fig. 2. System model

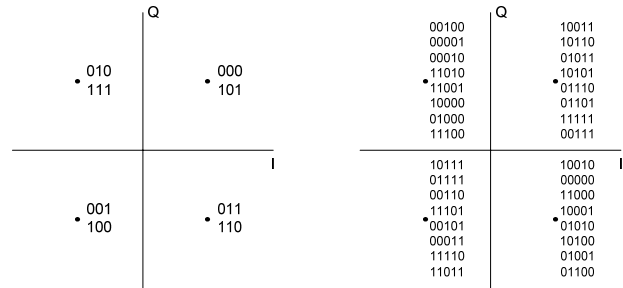


Fig. 1. Extended mapping 4-QAM $l_{map}=3$ and $l_{map}=5$, optimized with *a priori* information

inter-symbol interference. However, combining the technique presented in this paper with the turbo equalization framework [10]-[13] is rather straightforward.

It is assumed that transmission chain is properly normalized so that the received symbol energy-to-noise power spectral density ratio $E_s/N_0 = 1/\sigma_n^2$; with this normalization, we can properly delete the channel complex gain term from the mathematical expression of the channel. The discrete time description of the received signal $y(k)$ is then expressed by

$$y(k) = x(k) + n(k), \quad (1)$$

where, with k being the symbol timing index, $x(k)$ is the transmitted modulated signal with unit power, and $n(k)$ the zero mean complex AWGN component with variance σ_n^2 (i.e., $\langle |x(k)|^2 \rangle = 1$, $\langle n(k) \rangle = 0$, and $\langle |n(k)|^2 \rangle = \sigma_n^2$).

Receiver

At the receiver side, the iterative processing is invoked, where extrinsic information is exchanged between the demapper and decoder. The extrinsic LLRs calculated by soft input soft output (SISO) decoding/demapping are fed back and used for demapping/decoding as *a priori* LLR; the iteration continues until no more relevant gains in extrinsic MI can be achieved [14]; when such convergence point is reached, binary decisions are made on the information bits based on the *a posteriori* LLR at the variable node. Therefore,

the larger the MI at the convergence point, the lower the BER floor, which is depending on the matching between the encoder and the mapping rule.

The demapper calculates from the received signal point $y(k)$, corrupted by AWGN, the extrinsic LLR of the μ^{th} bit in the symbol transmitted at the k^{th} symbol timing, by

$$L_e[b_\mu(k)] = \ln \frac{\sum_{s \in S_0} \exp\left\{-\frac{|y-s|^2}{\sigma_N^2}\right\} \prod_{v=1, v \neq \mu}^{l_{\text{map}}} \exp(-b_v(s) L_a(b_v(s)))}{\sum_{s \in S_1} \exp\left\{-\frac{|y-s|^2}{\sigma_N^2}\right\} \prod_{v=1, v \neq \mu}^{l_{\text{map}}} \exp(-b_v(s) L_a(b_v(s)))} \quad (2)$$

where s denotes a signal point in the constellation, $S_0(S_1)$ indicates the set of the labels having the μ^{th} bit being 0(1), and $L_a(b_v(s))$ the a priori LLR fed back from the decoder corresponding to the v^{th} position in the label allocated to the signal point s .

A. Repetition Coded EM BICM-ID

Fig. 3 (A) shows a block diagram of the decoder with the structure (A). The d_v bits constituting one segment, output from the de-interleaver, are connected to one variable node, where the a priori LLRs to be updated is calculated by summing up the (d_v-1) LLRs, as

$$L_{e,j} = \sum_{i=1, i \neq j}^{d_v} L_{a,i} \quad (3)$$

to produce the extrinsic LLR for the j^{th} bit in the segment. This process is performed for all the other bits in the same segment as well as for all the other segments independently in the frame. Finally, the updated extrinsic LLRs are interleaved, and fed back to the demapper. With the a priori LLRs provided in the form of the decoder's output extrinsic LLR, demapper again performs the processing of Eq. (2) to update the demapper output extrinsic LLRs. This process is repeated. Obviously, the rate of this code is $1/d_v$, and the spectrum efficiency is l_{map}/d_v bits per symbol.

B. Repetition Coded EM BICM-ID with Check Nodes

Fig. 3 (B) shows a block diagram of the decoder with the structure (B). The d_v bits constituting one segment, output from the de-interleaver, are connected to a variable node, and d_c variable nodes are further segmented and connected to a check node decoder; those demapper output bits in one segment, connected to the same variable node decoder, are not overlapping with other segments, and so is the case of the variable node segmentation. Therefore, no iterations in the decoder are required. The extrinsic LLR update for a bit at the check node is exactly the same as the check node operation in the LDPC codes, as

$$L_{e,\text{cnd},k} = \sum_{i=1, i \neq k}^{d_c} \boxplus L_{a,\text{cnd},i} \quad (4)$$

where $\sum \boxplus$ indicates the box-sum operator [15].

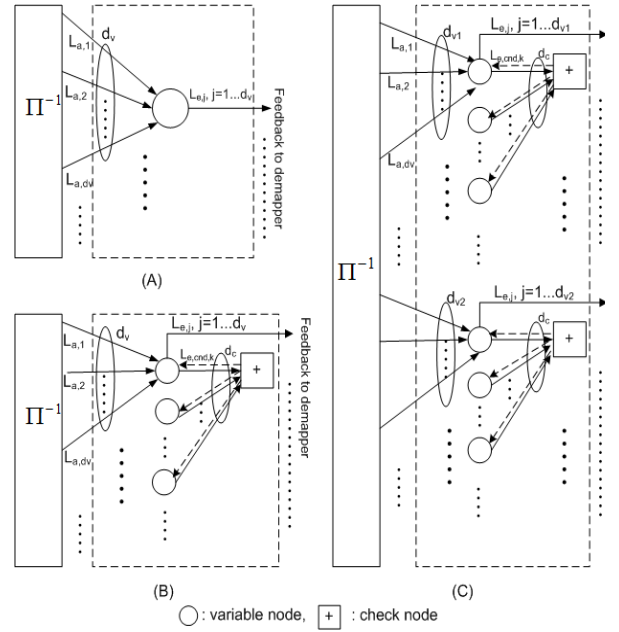


Fig. 3. Structures of the decoders

The extrinsic LLR, calculated by the check node decoder, is fed back to its connected variable node decoder, where it is further combined with (d_v-1) a priori LLRs forwarded from the demapper via the de-interleaver, as

$$L_{e,j} = L_{e,\text{cnd},k} + \sum_{i=1, i \neq j}^{d_v} L_{a,i} \quad (5)$$

This process is performed for the other bits in the same segment, and also for all the other segments independently in the same transmission block. Finally, the updated extrinsic LLRs obtained at the each variable node are interleaved, and fed back to the demapper. With the a priori LLRs provided in the form of decoder's output extrinsic LLR, demapper again performs the processing of (2) to update the demapper output extrinsic LLRs. This process is repeated. The rate of this code is $(d_c-1)/(d_c \cdot d_v)$, and the spectrum efficiency is $l_{\text{map}} \cdot (d_c-1)/(d_c \cdot d_v)$ bits per symbol.

C. Irregular Degree Allocations

Fig. 3 (C) shows a block diagram of the decoder with the structure (C). The segment-wise structure of (C) is exactly the same as that of (B), but the variable node degrees $d_{v,i}$ may have different values segment-by-segment. The equations for the LLR update are also the same as that for (B), but the calculations have to reflect different values $d_{v,i}$ of the variable node degrees. The rate of the code is $(d_c-1)/(d_c \cdot \sum a_i \cdot d_{v,i})$, and the spectrum efficiency is $l_{\text{map}} \cdot (d_c-1)/(d_c \cdot \sum a_i \cdot d_{v,i})$ bits per symbol, where a_i represents the ratio of variable nodes having degree $d_{v,i}$ in a block.

III. DESIGN BASED ON EXIT CHART

Fig. 4 shows EXIT curves for 4-QAM demappers with l_{map} as a parameter, where the labels were determined so that the EXIT curve has the largest decay. Also, the EXIT curve

with Gray mapping is presented in the figure. It is found that the Gray mapping has a completely flat EXIT curve, while anti-Gray mapping has a decay. With $l_{map}=m$, since mapping rule does not change the constellation constraint capacity (CCC), the area under the EXIT curves has to stay almost the same so far as the same modulation is used ($m=constant$). With EM, the left most point is further decreased, and the right most point increased. However, the area under the curve is smaller than without EM, because even with $l_{map}>m$, still the spectrum efficiency of the modulation itself stays the same ($=m$).

(A) Repetition Code

With the Gaussian assumption for the LLR distribution, the EXIT function of the repetition code is given by

$$I_{e,v} = J\left(\sqrt{(d_v - 1) \cdot J^{-1}(I_{a,v})^2}\right). \quad (6)$$

where $I_{a,v}$ is the variable node input a priori MI and $I_{e,v}$ is its output extrinsic MI. $J()$ and $J^{-1}()$ are the functions that convert the square-root variance σ of LLR to its corresponding MI, and its inverse, respectively [2]. Obviously, (6) is corresponding to (3) for LLR update, with which $I_{a,v} = I_{e,dem}$ with $I_{e,dem}$ being the demapper output extrinsic MI.

(B) Repetition Code with Check Node

The check node EXIT function can be approximated by [16]

$$I_{e,cnd} = 1 - J\left(\sqrt{d_c - 1} \cdot J^{-1}(1 - I_{a,cnd})\right), \quad (7)$$

where

$$I_{a,cnd} = J\left(\sqrt{d_v \cdot J^{-1}(I_{a,dec})^2}\right). \quad (8)$$

The EXIT function of the whole decoder comprised of the variable and check node decoders can be calculated by combining Eq. (6) and (7) [4], as

$$I_{e,dec} = J\left(\sqrt{(d_v - 1) \cdot J^{-1}(I_{a,dec})^2 + J^{-1}(I_{e,cnd})^2}\right), \quad (9)$$

with $I_{a,dec} = I_{e,dem}$.

(C) Irregular Repetition Code with Check Node

The EXIT function of the whole decoder with the structure (C) can be obtained by weighting the segment-wise EXIT functions given by (9), by a_i corresponding to their distributions, as

$$I_{e,dec} = \frac{\sum_i a_i \cdot d_{v,i} \cdot J\left(\sqrt{(d_{v,i} - 1) \cdot J^{-1}(I_{a,dec})^2 + J^{-1}(I_{e,cnd})^2}\right)}{\sum_i a_i \cdot d_{v,i}}. \quad (10)$$

Hence, the shape of the combined code EXIT function can be flexibly controlled so that better matching between demapper and decoder curves can be achieved.

Fig. 5 shows EXIT functions of the decoders (A), (B), and (C) for the degree allocations presented in the box in the

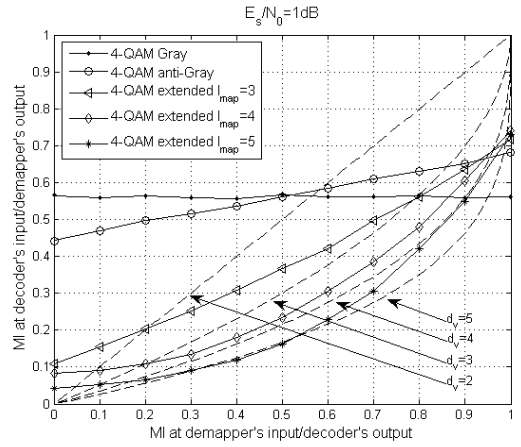


Fig. 4. EXIT chart: 4-QAM extended mapping $l_{map}=2, \dots, 5$ combined with repetition code $d_v=2, \dots, 5$

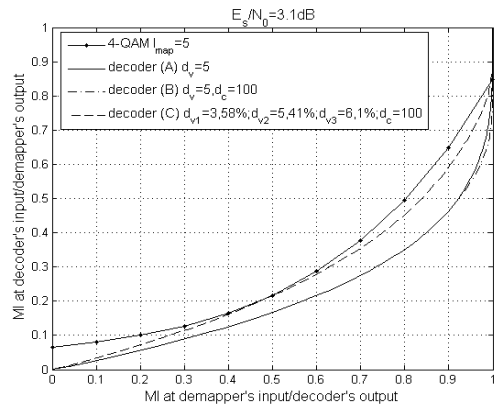
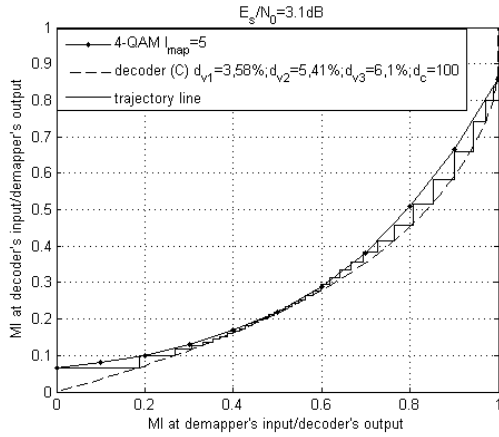


Fig. 5. EXIT chart: 4-QAM EM $l_{map}=5$ combined with 3 kinds of repetition code due to structures (A), (B) and (C)

figure, where the code rates are 1, 0.99, and 1.29, respectively. It is found by comparing Figs. 4 and 5 that the EXIT functions of EM and the repetition code with check nodes can be well matched by changing the degree allocations a_i , by which we can expect sharp turbo cliff to happen at their corresponding threshold E_s/N_0 , even though the structures of (A), (B), and (C) are very simple and easy to implement.

IV. NUMERICAL RESULTS

A series of simulations was conducted with enough number of bits transmitted to verify the advantageous characteristics of the techniques proposed. As described above, the EM technique does not change the CCC, and therefore the capability of achieving near capacity performance even with the very simple structure, which is the most significant advantage of the proposed technique, is mostly in a low E_s/N_0 value range (or equivalently low rate code case). If the CCC is much lower than the Gaussian capacity at the system's operation E_s/N_0 point, no significant merit can be expected, and obviously larger modulation format should be used in such case.


 Fig. 6. EXIT chart: 4-QAM EM $l_{map}=5$ combined with repetition code $d_v=5$

With irregular degree allocation of (C), as shown by Fig. 6, the trajectory sneaks through the narrowly tiny open tunnel, and finally reaches the intersection point. With the structure (A), the adjustable parameters are only l_{map} and d_v . We therefore have introduced the check node in (B), and furthermore, irregular degree allocation in (C), both to (A). Two BER curves with $l_{map}=5$ EM 4-QAM with structure (C) are shown in Fig. 7, together with $l_{map}=5$ EM 4-QAM and $d_v=5$ to demonstrate the design flexibility. The degree allocations are shown under the figure caption. In fact, the allocation parameters for the rate 1.29 code, of which BER is indicated by “o” in the figure, are exactly the same as that shown in Fig. 5 for (C) with $E_s/N_0=3.1dB$. It is found that the EXIT and BER curves, presented in Figs. 5 and 8, are exactly consistent with each other (at $E_s/N_0=3.1dB$, EXIT tunnel opens and BER turbo cliff happens). These observations indicate that with the design flexibility made available by introducing the irregular structure, we can flexibly control the threshold E_s/N_0 and the error floor.

V. MODULATION DOPING

The technique described above is flexible in making good matching between the demapper and decoder EXIT curves, but is not suitable in designing BICM-ID with very low rate codes that achieve the BER pinch-off at very low SNR. This is because the LM point of the demapper EXIT function, corresponding to decoder feedback mutual information $MI=0$, is already very low, and there is not enough room to further lower the EXIT LM point while avoiding the intersection between the demapper and decoder curves. To solve this problem, this paper applies the idea of mixing modulation symbols having different labeling rules for the extended anti-Gray mapping and the standard Gray mapping, at a certain ratio. Since the demapper EXIT function with Gray mapping is relatively flat, its LM point has the higher MI value than EM mapping rules. Therefore, with modulation doping the demapper EXIT LM point is lifted up from that without doping, and the amount depends on the doping ratio.

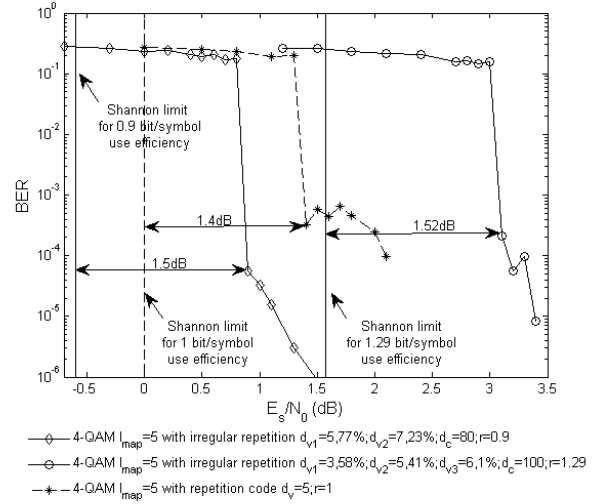


Fig. 7. BER curves optimized by irregular repetition code

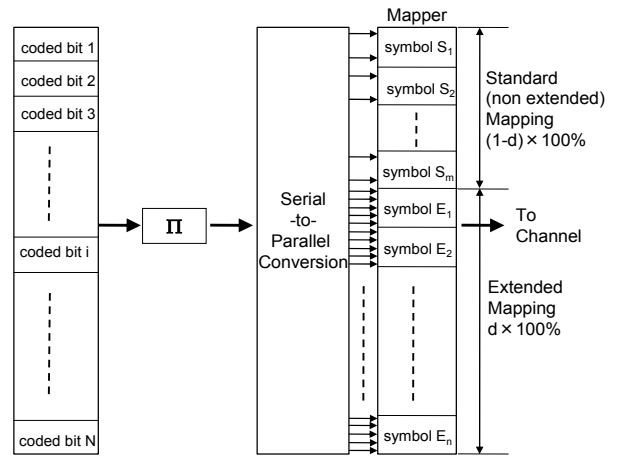


Fig. 8. Block Diagram of Modulation Doping

With EM, more than one labels having different bit patterns in the segment are mapped on to each constellation point ($l_{map} > m$), while without mapping extension, such as standard Gray mapping, only one bit pattern in the segment is mapped on to each constellation point ($l_{map} = m$). Fig. 8 shows a block diagram of the modulation doping system. The spectrum efficiency of the system then becomes $\eta = \{d \cdot m + (1-d) \cdot l_{map}\} R_c$, where d and $(1-d)$ are the ratios of the symbols with doped (standard) mapping and extended mappings, respectively, in a transmission frame. R_c is the code rate, given by $R_c = (d_c - 1) / (d_c \cdot \sum a_i \cdot d_{v,i})$, where a_i is the ratio of the variable node having degree $d_{v,i}$. It has been assumed that the interleaver is long and random enough so that the repetition encoder output bits are uniformly spread over the whole transmission frame.

Fig. 9 shows EXIT curves of some representative design examples using EM 4QAM with $l_{map}=5$: Their corresponding

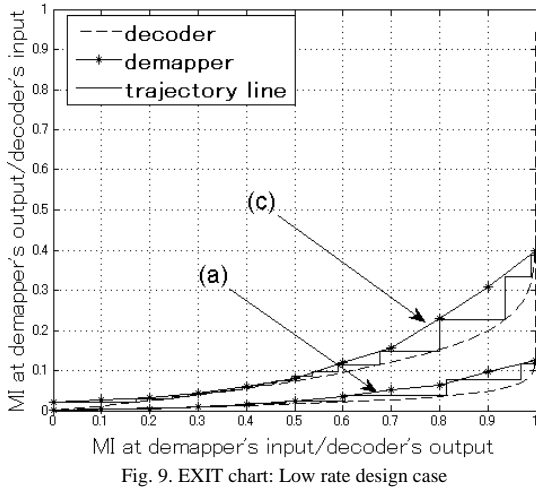


Table 1. Degree allocations of the codes shown in Fig. 9

	E_s/N_0 (dB)	$d_{v1},$ a_1	$d_{v2},$ a_2	$d_{v3},$ a_3	d_c	d	bits/symbol
(a)	-9.2	30, 0.5	40, 0.5	/	3	0.03	0.0935
(b)	-6.3	18, 0.5	20, 0.5	/	5	0.04	0.2055
(c)	-3.4	9, 0.5	10, 0.05	12, 0.45	8	0.05	0.4081

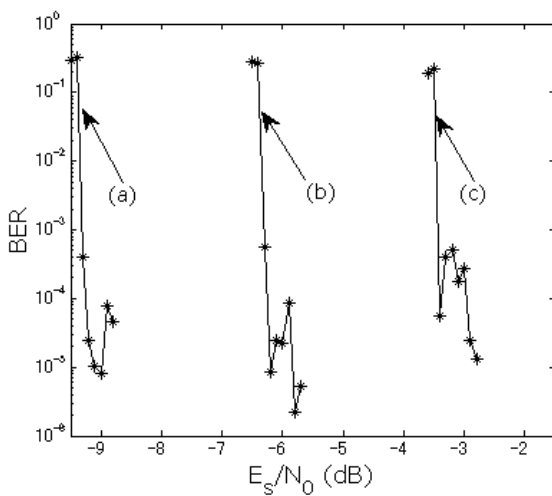


Fig. 10. BER Performance for the code design of Table 1

code design parameters and their threshold E_s/N_0 values are summarized in Table 1. A series of chain simulations for the code designs shown in Table 1 has been conducted to verify the advantageous characteristics of the techniques proposed

in this paper. The results of the chain simulations conducted to obtain the trajectories for the code parameters shown in Table 1 are also shown in Fig. 9. It should be emphasized that the trajectory and EXIT curve are exactly consistent each other. Fig. 10 shows the BER performance for the code design of Table 1. It is found that the BER pinch-off happens exactly at the values indicated by the EXIT chart analysis.

VI. CONCLUSIONS

In this paper, a new BICM-ID technique that makes efficient use of a mixture of standard Gray and anti-Gray extended mapping, combined with irregular degree-allocated repetition code with check nodes has been proposed. Mathematical details as well as EXIT analysis and BER simulation results have been provided. It has been shown that with the proposed technique, very low rate code that can achieve a turbo cliff at a very low E_s/N_0 value range can well be designed, and that the EXIT chart analysis and chain simulation results are exactly consistent with each other.

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